## Sparsity — tutorial 3

Measuring sparsity

**Problem 2.** Suppose  $\mathcal{C}$  is a class of bounded expansion. Prove that for every  $r \in \mathbb{N}$  there exists a constant  $c_r$  such that the following holds. For every graph  $G \in \mathcal{C}$  and every its vertex subset  $A \subseteq V(G)$ , there exists a vertex subset  $B \supseteq A$  with the following properties:

- $|B| \leq c_r |A|$ , and
- for every pair of vertices  $u, v \in A$ , if  $\operatorname{dist}_G(u, v) \leq r$  then  $\operatorname{dist}_{G[B]}(u, v) = \operatorname{dist}_G(u, v)$ .

**Solution.** We will use the following two facts; the first one was proved in the first tutorial, the second one was given in the second.

Fact 1 (Closure lemma). Suppose  $\mathcal{C}$  is a class of bounded expansion. Then for every  $r \in \mathbb{N}$  there exists a constant  $d_r$  such that the following holds. For every graph  $G \in \mathcal{C}$  and every subset A of its vertices, there exists a vertex subset  $B \supseteq A$  such that  $|B| \leqslant d_r |A|$  and for every vertex  $u \in V(G) - B$ , at most  $d_r$  vertices of B can be reached from u by a path of length at most r whose internal vertices do not belong to B.

**Fact 2** (Stability under lexproduct). If G is a graph and  $r, c \in \mathbb{N}$ , then

$$\nabla_r(G \bullet K_c) \leq 2c^2(r+1)^2 \nabla_r(G) + c.$$

We now present a solution to Problem 1. First, we apply Lemma 1 to the set A, yielding a set A' with the asserted properties:  $|A'| \leq d_r |A|$  and for every vertex  $u \in V(G) - A'$ , at most  $d_r$  vertices of A' can be reached from u by a path of length at most r whose internal vertices do not belong to A'. Next, for each pair of distinct vertices  $u, v \in A'$ , select an arbitrary path  $P_{u,v}$  that connects u and v, and whose internal vertices do not belong to A', and which is the shortest among the paths satisfying these properties; in case there is no such path, put  $P_{u,v} = \emptyset$ . Then define B to be A' plus the vertex sets of all paths  $P_{u,v}$  that have length at most r.

We first claim that B has indeed the required property of preserving distances up to r. More precisely, take any distinct  $u, v \in A$  with  $\operatorname{dist}_G(u, v) \leqslant r$ . Let R be a shortest path between u and v in G, and let  $a_1, a_2, \ldots, a_q$  be consecutive vertices of A' visited on R, where  $u = a_1$  and  $v = a_q$ . For each  $i = 1, 2, \ldots, q-1$ , let  $R_i$  be the segment of R between  $a_i$  and  $a_{i+1}$ . Then the existence of  $R_i$  certifies that some path of length at most  $|R_i|$  between  $a_i$  and  $a_{i+1}$  was added when constructing B from A', and hence  $\operatorname{dist}_{G[B]}(a_i, a_{i+1}) \leqslant |R_i|$ . Consequently, by the triangle inequality we infer that

$$\operatorname{dist}_{G[B]}(u,v) \leqslant \sum_{i=1}^{q-1} \operatorname{dist}_{G[B]}(a_i, a_{i+1}) \leqslant \sum_{i=1}^{q-1} |R_i| = |R| = \operatorname{dist}_G(u, v).$$

However, the opposite inequality  $\operatorname{dist}_{G[B]}(u,v) \ge \operatorname{dist}_{G}(u,v)$  follows directly from the fact that G[B] is an induced subgraph of G. Hence indeed  $\operatorname{dist}_{G}(u,v) = \operatorname{dist}_{G[B]}(u,v)$ .

We are left with showing that  $|B| \leq c_r |A|$  for some constant  $c_r$ . First, we have  $|A'| \leq d_r |A|$ , so we only need to upper bound the ratio  $\frac{|B|}{|A'|}$ . Let H be a graph on vertex set A', where  $uv \in E(A')$  if and only if  $P_{u,v}$  exists and has length at most r, and hence its vertex set was added in the costruction of B. Clearly  $|B| \leq |A'| + (r-1) \cdot |E(H)|$ , so it suffices to prove an upper bound on |E(H)|.

Take any  $w \in B - A'$ , and consider for how many pairs  $\{u, v\}$  it can hold that  $w \in P_{u,v}$ . If  $\{u, v\}$  is such a pair, then in particular both u and v can be reached from w by a path of length at most r that internally avoids A'. However, we know that the number of such vertices is at most  $d_r$ , so the number of such pairs  $\{u, v\}$  is at most  $\tau = \binom{d_r}{2}$ . Consequently, we observe that graph H is an (r-1)-shallow minor (actually

even an  $\lceil (r-1)/2 \rceil$ -shallow topological minor) of  $G \bullet K_{\tau}$ : when each vertex  $w \in B-A'$  is replaced with  $\tau$  copies, then we can realize all the paths  $P_{u,v}$  in  $G \bullet K_{\tau}$  so that they are internally vertex-disjoint. Now we know by Fact 1 that  $\nabla_{r-1}(G \bullet K_{\tau})$  is bounded by a function of  $\nabla_{r-1}(G)$  and  $\tau$ . Both  $\nabla_{r-1}(G)$  and  $\tau$  are bounded by constants, namely by  $\nabla_{r-1}(\mathcal{C})$  and  $\binom{d_r}{2}$  respectively, and hence so does  $\nabla_{r-1}(G \bullet K_{\tau})$ . As  $|E(H)| \leqslant \nabla_{r-1}(G \bullet K_{\tau}) \cdot |A'|, |B| \leqslant |A'| + (r-1)|E(H)|, \text{ and } |A'| \leqslant c_r|A|, \text{ we are done.}$