

Problem set II, for October 14th

Short introduction to categories.

It is assumed that you know basic definitions of this theory. There are many sources for getting them; here are two of them:

- Adamek, Herrlich, Stecker, *Abstract and concrete categories*
- Freyd, *Abelian Categories*

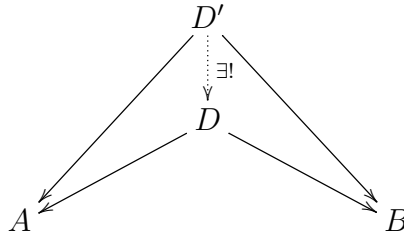
The categories of sets and algebraic structures (groups, abelian groups, commutative rings, vector spaces over a fixed field k), and topological space will be denoted by, respectively, \mathcal{Set} , \mathcal{Gr} , \mathcal{Ab} , \mathcal{Ring} , \mathcal{Vect}_k , \mathcal{Top} .

1. Define the notion of an isomorphism and of the inverse of a morphism in categorical terms.
2. An object A in a category \mathcal{C} is called initial if it has exactly one morphism into every object of \mathcal{C} . Similarly we define the terminal object of \mathcal{C} . Prove: if an initial (terminal) object exists then it is unique up to isomorphism. Check whether the categories listed above have the initial (terminal) object.
3. Prove that abelianization is a covariant functor $Ab : \mathcal{Gr} \rightarrow \mathcal{Ab}$. Is this a full functor? Is this a faithful functor?
4. Let (S, \leq) be a set with partial order. Show that (S, \leq) determines a category with objects being elements from S and morphisms determined by the order.
5. Let \mathcal{C} be a category with a distinguished object A . Show that mappings $B \rightarrow \text{Mor}_{\mathcal{C}}(A, B)$ and

$$\text{Mor}_{\mathcal{C}}(B, C) \ni f \longrightarrow (\text{Mor}_{\mathcal{C}}(A, B) \ni g \rightarrow f \circ g \in \text{Mor}_{\mathcal{C}}(A, C))$$

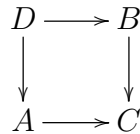
define covariant functor from \mathcal{C} into \mathcal{Set} . We will denote it by $\mathfrak{h}_A : B \rightarrow \text{Mor}_{\mathcal{C}}(A, B)$. Similarly we define a covariant functor $\mathfrak{h}^A : B \rightarrow \text{Mor}_{\mathcal{C}}(B, A)$. Prove that any morphism $A' \rightarrow A$ determines natural transformations of functors $\mathfrak{h}_A \rightarrow \mathfrak{h}_{A'}$ and $\mathfrak{h}^{A'} \rightarrow \mathfrak{h}^A$.

6. Yoneda lemma. For a given category \mathcal{C} we define its functor category $\widehat{\mathcal{C}}$ whose objects are contravariant functors $\mathcal{C} \rightarrow \mathcal{Set}$ and morphisms are natural transformations of such functors. Let us consider a functor $\widehat{\Phi} : \mathcal{C} \rightarrow \widehat{\mathcal{C}}$ such that $\widehat{\Phi}_{\text{Obj}}(A) = \mathfrak{h}^A$ while on morphisms $\widehat{\Phi}$ is as defined in the previous exercise. Show that $\widehat{\Phi}$ is faithful and full.
7. A contravariant functor $\mathcal{C} \rightarrow \mathcal{Set}$ is called representable by an object A in \mathcal{C} if it is naturally isomorphic to the functor \mathfrak{h}^A . Show that the forgetful functor $\mathcal{Ab} \rightarrow \mathcal{Set}$ is naturally isomorphic to $\mathfrak{h}_{\mathbb{Z}}$.
8. Let $\text{vect}_k^{<\infty}$ denote category in which objects are linear spaces k^n (one object for every $n \geq 0$) while $\text{Mor}_{\text{vect}}(k^n, k^m)$ are matrices $n \times m$ with multiplication as composition. Show that the embedding of $\text{vect}_k^{<\infty}$ into $\mathcal{Vect}_k^{<\infty}$ is equivalence of categories.
9. Categorical product of two objects $A, B \in \text{Obj}_{\mathcal{C}}$ is an object D with two morphisms $A \leftarrow D \rightarrow B$ such that for any object D' with morphisms $A \leftarrow D' \rightarrow B$ there exists exactly one morphism from D' to D which makes the following diagram commutative



Usually D is denoted by $A \times B$. Co-product is defined similarly by reversing the arrows. Prove that the product (co-product) is unique up to isomorphism. Prove that $A \times B$ is naturally isomorphic to $B \times A$. Describe products and co-products in \mathcal{Set} , \mathcal{Top} , \mathcal{Ab} , \mathcal{Gr} , \mathcal{Ring} ; note that they do not have to exist.

10. In category \mathcal{C} we consider three objects with two morphisms $A \rightarrow C \leftarrow B$. The fiber product of A and B over C is an object D with morphisms $A \leftarrow D \rightarrow B$ which make this diagram commutative



Moreover if D' satisfies the above condition formulated for D then there exists a unique morphism $D' \rightarrow D$ which makes the following diagram commutative

$$\begin{array}{ccccc}
 D' & & & & \\
 & \searrow & & \searrow & \\
 & & D & \longrightarrow & B \\
 & \searrow & \downarrow & & \downarrow \\
 & & A & \longrightarrow & C
 \end{array}$$

We denote such product (if it exists) by $A \times_C B$. The fiber co-product is defined similarly by reversing the arrows. Show the uniqueness (up to isomorphism) of the fiber product and co-product. Contemplate the fiber products and co-products in \mathbf{Set} , \mathbf{Top} , \mathbf{Ab} . Prove that if \mathcal{C} has the terminal object Z and admits the product $A \times B$ then the fiber product $A \times_Z B$ is isomorphic to $A \times B$.

11. Composition of fiber products is a fiber product. Suppose that the following arrows are defined

$$\begin{array}{ccccc}
 E \times_A (A \times_C B) & \longrightarrow & A \times_C B & \longrightarrow & B \\
 \downarrow & & \downarrow & & \downarrow \\
 E & \longrightarrow & A & \longrightarrow & C
 \end{array}$$

Prove that $E \times_A (A \times_C B)$ is isomorphic to $E \times_C B$ where $E \rightarrow C$ is the composition $E \rightarrow A \rightarrow C$.

12. Fiber product base change. Suppose that there exists a fiber product for $A \rightarrow C \leftarrow B$. Prove that for every $C \rightarrow E$ we have a natural morphism $A \times_C B \rightarrow A \times_E B$ and the diagram

$$\begin{array}{ccc}
 A \times_C B & \longrightarrow & A \times_E B \\
 \downarrow & & \downarrow \\
 C & \longrightarrow & C \times_E C
 \end{array}$$

where $C \rightarrow C \times_E C$ is the diagonal which means the morphism coming from two identities $C \rightarrow C \leftarrow C$. Prove that the above diagram makes the fiber product.