## Proportionality Guarantees in Elections with Interdependent Issues Markus Brill 14, Evangelos Markakis 23, Georgios (Yorgos) Papasotiropoulos 2, Jannik Peters 4

<sup>1</sup>University of Warwick, UK. <sup>2</sup>Athens University of Economics and Business, Greece. <sup>3</sup>Input Output Global (IOG). <sup>4</sup>TU Berlin, Germany.



...some voters may be able to express their opinion using approval ballots...

(Pizza? Burger? Yummy! I do like donuts and choco-bars! Don't you see that I am too young for alcohol?



... but others may not; mainly due to dependencies between issues!

(I'd like a burger but only if paired with a beer. I don't have a strong opinion when it comes to desserts. I can't have both crab and chocolate!

[BL16]: framework for expressing dependencies in the context of AV

[SG22]: proportionality criterion for binary and independent issues & powerful guarantees for.PAV and MES

> For every sufficiently large group of voters, independent of cohesiveness requirements.

Any group of voters that makes up an x-fraction of the electorate, should be able to decide on an x-fraction of the issues.



Wow! Combining [BL16] with [SG22] might lead to provable proportionality guarantees for elections with interdependent issues! Consider an election with **n** voters and **m** issues, each of domain **d**, s.t. voter  $v_i$  casts dependency graph  $G_i$  and conditional approval ballots.

G := undirected variant of  $\bigcup G_i$ 

:= maximum degree in G

 $u_i(w) :=$  satisfaction of voter  $v_i$ , under outcome  $w \in d^m$ 

For a set of voters V,  $\mathbf{r}_i := \#$  issues: every  $v_i \in V$  approves  $\geq 1$  alternative.

A rule R is *a-proportional, a*  $\in$  [0,1], if for every set of voters V, there exists a voter v<sub>i</sub>  $\in$  V s.t. if w: winning outcome under R,

then  $u_i(w) > \alpha r_v \frac{|V|}{n} - 1$ , for any instance.

Conditional Proportional Approval Voting (CPAV)

winning outcome :=  $argmax_{w \in d^m} \sum_{V_i \in N} \sum_{k=1}^{m} 1/k^2$ 

## Conditional Method of Equal Shares (CMES)

- ★ assign a budget of m to every voter
- ★ for every yet unfixed issue  $I_j$  and for every possible (sub)outcome w for some k issues in the closed neighborhood of  $I_j$  in G:
- ★ S(w) = voters with >0 budget, satisfied w.r.t. I<sub>j</sub> under w.
- ★ p(w) = price s.t. if every voter in S(w) paid p(w) or all the money she has left, then voters from S(w) would altogether pay nk.
- ★ if no purchase can be made, fix remaining issues arbitrarily, otherwise, select w that minimizes p(w) and reduce voters' budget accordingly. Then, repeat from Step 2

observation 1.

Any "reasonably fair" rule,

cannot be  $\alpha$ -proportional for  $\alpha > 1/a^{\Delta+1}$ .

Assumption 1: Voters that can be satisfied w.r.t. to an issue, cannot be satisfied w.r.t. to other 'nearby' issues in G.

Under Assumption 1, cPAV is  $\alpha$ -proportional, for  $a = \frac{1}{(\alpha + \Delta^{\alpha})} a^{\Delta + 1}$ .  $\Delta = 0: \alpha = \frac{1}{\alpha} \text{ strict generalization}$ of the  $\frac{1}{2}$  factor (binary issues) [SG22] (tight according to Observation 1) For  $\Delta > 0$ , cMES cannot be  $\alpha$ -proportional, for any  $\alpha$ , even under Assumption 1.

Assumption 2: For every voter  $v_i$ , for every issue  $I_j$  and for every combination of alternatives for issues in the in-neighborhood of  $I_j$  in  $G_i$ , there is an alternative of  $I_i$  that satisfies  $v_i$ .



[BL16]: N. Barrot & J. Lang. Conditional and Sequential Approval Voting on Combinatorial Domains. In International Joint Conference on AI (IJCAI), 2016. [SG22]: P. Skowron & A. Gorecki. Proportional Public Decisions. In AAAI Conference on AI (AAAI), 2022.